

The regular completion of **Cat** and its relation to lazy categories

Robert Paré
(Dalhousie University)

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Introduction

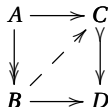
- The theory of lazy categories aims to fix quotients in **Cat** by embedding it in a nice topos (Verity)
 - What's wrong with quotients in **Cat**?
 - Perhaps the main thing is that they are not stable under pullback
 - But they just seem wrong – unruly and just plain nasty
 - Exercise: Describe the monoid of endomorphisms $X(2,2)$ for the coequalizer

$$\mathbf{Set}_0 \begin{array}{c} \xrightarrow{2 \times ()} \\ \xrightarrow{3 \times ()} \end{array} \mathbf{Set}_0 \longrightarrow\!\!\!\twoheadrightarrow X$$

- Another way of “fixing quotients” is the regular completion of Carboni
- We compare the two

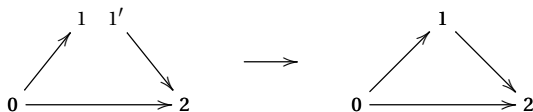
Regular categories

- **Regular categories** are categories with a good factorization system
 - finite limits
 - every morphism factors as a regular epi followed by a mono
 - regular epis are stable under pullback
- A regular epimorphism is a coequalizer of some pair of morphisms
- They are orthogonal to monomorphisms



Examples

- Algebraic categories
- Any topos
- But not **Cat**, ex:



does not factor

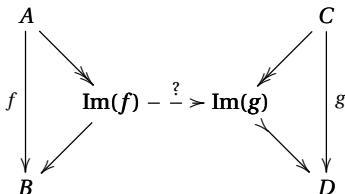
Regular completion

- A **regular functor** is a functor that preserves finite limits (i.e. is **left exact**) and regular epimorphisms
- We have a forgetful 2-functor

$$U: \mathcal{R}eg \longrightarrow \mathcal{L}ex$$

and we want a left (2-)adjoint for it

- Idea: Formally add images



The construction (Carboni *et al.*)

$\mathbf{A}$ with finite limits

$\text{Reg}(\mathbf{A})$

- objects: morphisms of $\mathbf{A}$
- morphism: $h: (f) \twoheadrightarrow (g)$

$$\begin{array}{ccc} A & \overset{\bar{h}}{\dashrightarrow} & C \\ f \downarrow & \searrow h & \downarrow g \\ B & & D \end{array}$$

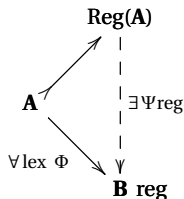
h lifts to some $\bar{h}$

coequalizes the kernel pair of f

- Terminal object is $1: 1 \twoheadrightarrow 1$
- Pulbacks = pullbacks of h 's
- Factorization

$$\begin{array}{ccccc} A & \overset{1}{\dashrightarrow} & A & \overset{\bar{h}}{\dashrightarrow} & C \\ f \downarrow & \searrow h & \downarrow h & \searrow h & \downarrow g \\ B & & D & & D \end{array}$$

Universal property



Theorem

Reg is a (pseudo) left adjoint to $U: \mathcal{R}eg \longrightarrow \mathcal{L}ex$, and U is 2-monadic.

- $\Psi(f) = \text{Im}(\Phi(f))$.
- We can apply Reg to $\mathbf{Cat}$ (size considerations notwithstanding) and get an extension of $\mathbf{Cat}$

$$\mathbf{Cat} \longrightarrow \text{Reg}(\mathbf{Cat})$$

where images (and relations) work well.

- But there's another way – the lazy way!

The topos hull of **Cat**

- Embed **Cat** into the most fitting Grothendieck topos – sheaves on the canonical topology
- The covering families $\langle F_i: \mathbf{A}_i \longrightarrow \mathbf{B} \rangle$ are those for which the induced functor

$$F: \Sigma \mathbf{A}_i \twoheadrightarrow \mathbf{B}$$

is onto on composable pairs

- The $F: \mathbf{A} \twoheadrightarrow \mathbf{B}$ that are onto on pairs are precisely the stable regular epis in **Cat**: we call them *covers*
- $\text{Sh}(\mathbf{Cat}, \text{Cov}) \simeq \mathbf{Set}^{\Delta_2^{op}}$

Lazy categories

Definition

A *lazy category* is a diagram of sets and functions

$$A_2 \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\lambda} \\ \xrightarrow{\quad} \\ \xleftarrow{\rho} \\ \xrightarrow{\quad} \end{array} A_1 \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\text{id}} \\ \xrightarrow{\quad} \end{array} A_0$$

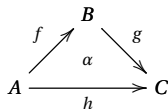
satisfying well-known equations.

A *lazy functor* is given by three functions commuting with the structure maps.

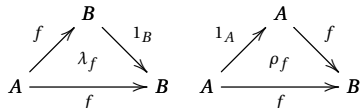
- We have a full embedding $\mathbf{Cat} \longrightarrow \mathbf{LCat}$
 - $\mathbf{LCat}$ is a Grothendieck topos, so it's like $\mathbf{Set}$
 - But it's also like $\mathbf{Cat}$

The device

- Introduce category like notation
 - Elements of A_0 are called *objects*, denoted $A, B, C, \dots$
 - Elements of A_1 are called *morphisms*, denoted $f: A \longrightarrow B, \dots$
 - Elements of A_2 are called *composers*, denoted



- Identities $1_A: A \longrightarrow A$ and



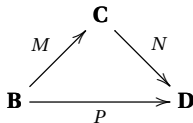
$$\lambda_{1_A} = \rho_{1_A} = \iota_A$$

- Functors preserve all of the above

Example

Let $\mathbf{A}$ be a category and define a lazy category $\text{Sub}(\mathbf{A})$ as follows

- Objects are subcategories $\mathbf{B} \rightrightarrows \mathbf{A}$
- Morphisms $M: \mathbf{B} \rightrightarrows \mathbf{C}$ are subsets $M \subseteq \text{Mor}(\mathbf{A})$ whose elements have domain in $\mathbf{B}$ and codomain in $\mathbf{C}$ and closed under composition by morphisms in $\mathbf{B}$ and $\mathbf{C}$ (ideal)
- Composers



contains $NM \subseteq P$

- For any category $\mathbf{X}$ we have a bijection

$$\frac{\mathbf{X} \rightrightarrows \text{Sub}(\mathbf{A}) \text{ in } \mathbf{LCat}}{\mathbf{B} \rightrightarrows \mathbf{A} \times \mathbf{X} \text{ in } \mathbf{Cat}}$$

Properties

Theorem

- (1) **Cat** $\longrightarrow$ **LCat** has a left adjoint **L**
- (2) **Cat** $\longrightarrow$ **LCat** preserves all limits
- (3) **Cat** $\longrightarrow$ **LCat** preserves coproducts and covers
- (4) **Cat** is an exponential ideal in **LCat**
- (5) Every lazy category has a cover by a projective category

Concerning (1) – For a lazy category **A**, **LA** is obtained by taking the free category on the morphisms of **A**, i.e. paths, and modding out by the congruence generated by identifying

$$A \xrightarrow{f} B \xrightarrow{g} C \quad \text{with} \quad A \xrightarrow{h} C$$

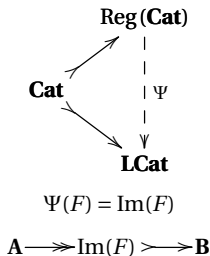
when there is a composer

$$\begin{array}{ccc} & B & \\ f \nearrow & & \searrow g \\ A & \xrightarrow{h} & C \\ & \alpha & \end{array}$$

Concerning (5) – The projectives in **LCat** are sums of 1, 2 and 3

$$\mathbf{LCat}(3, \mathbf{A}) \bullet 3 \longrightarrow \mathbf{A}$$

LCat vs Reg(Cat)



Theorem

- (1) Ψ has a left adjoint Λ
- (2) The induced monad $\Psi\Lambda$ is idempotent
- (3) The category of algebras (= fixed points) of $\Psi\Lambda$ is the full subcategory of **LCat** determined by those lazy categories embeddable in a category

Concerning (1)

$$\Lambda(A) = (\mathbf{LCat}(3, A) \bullet 3 \twoheadrightarrow A \rightarrow LA)$$

Partial categories

- A lazy subcategory of an actual category is called a *partial category*
- Every partial category is Ψ of an object of $\text{Reg}(\mathbf{Cat})$

$$\mathbf{LCat}(\mathbf{3}, \mathbf{A}) \bullet \mathbf{3} \longrightarrow \mathbf{A} \longrightarrow \mathbf{B}$$

- Every functor between partial categories is Ψ of a morphism in $\text{Reg}(\mathbf{Cat})$
- However Ψ is not full (it is faithful)

Properties of **ParCat**

Theorem

(1) *The inclusion **ParCat** $\hookrightarrow$ **LCat** has a left adjoint given by the image of $\eta\mathbf{A}: \mathbf{A} \rightarrow \mathbf{LA}$*

$$\mathbf{A} \twoheadrightarrow \text{Im}(\eta\mathbf{A}) \hookrightarrow \mathbf{LA}$$

(2) ***ParCat** is an exponential ideal in **LCat***

(3) ***ParCat** is regular*

(4) ***ParCat** is not locally cartesian closed*

Characteristics of partial categories

- Partial categories are *thin*, in fact *gaunt*

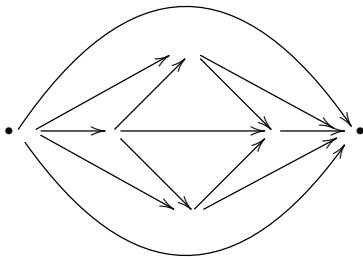
For any morphisms $A \xrightarrow{f} B \xrightarrow{g} C$ there is at most one composer

$$\begin{array}{ccc} & B & \\ f \nearrow & \alpha & \searrow g \\ A & \xrightarrow{h} & C \end{array}$$

- Partial categories are *weakly associative*: if $(fg)h$ and $f(gh)$ are both defined then they are equal. This holds for any two brackettings of an n -composite

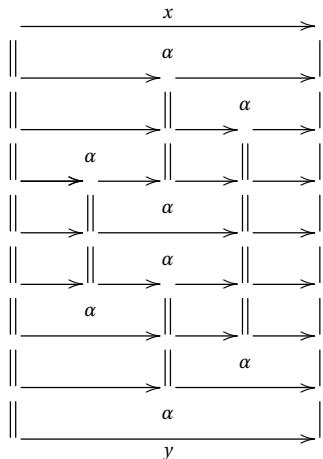
$$f_1 f_2 \dots f_n$$

- Not sufficient



Weird associativity?

- $\mathbf{A}$ is a partial category if and only if $\eta\mathbf{A}: \mathbf{A} \longrightarrow \mathbf{LA}$ is monic
- $x \sim y$



- **Question:** Can this be phrased as some kind of associativity?

Conclusion

- While $\text{Reg}(\mathbf{Cat})$ is an interesting category is not a good substitute for $\mathbf{Cat}$
 - Kills the good quotients
 - Unwieldy
 - Not even well-powered (Johnstone), so not locally presentable nor a topos
- Its image in $\mathbf{LCat}$ is $\mathbf{ParCat}$
 - Many interesting properties
 - Easy to work with
 - But it's not a topos (nor a quasi-topos)
 - Keep in mind ... but $\mathbf{LCat}$ rules!

Thank you,
and come again!