

Exploring double categories of relations

Hayato Nasu

Kyoto Univ.

this Sep.



Dalhousie Univ.

FMCS 2025

June 18 , 2025

Introduction

01 / 17

Fibrations

$\xrightarrow{\text{Bil}}$

Double categories

$$\text{Set}^{\subseteq} \xrightarrow{\text{cod}} \text{Set}$$

Rel

$$\text{Set}^{\rightarrow} \xrightarrow{\text{cod}} \text{Set}$$

Span

$$\mathcal{C}^{\rightrightarrows} \xrightarrow{\text{cod}} \mathcal{C}$$

$\mapsto$

$\text{Rel}(\mathcal{C})$

$$\text{Q-Fam} \xrightarrow{|\cdot|} \text{Set}$$

Q-Rel

$$M \hookrightarrow \mathcal{C}^{\rightarrow} \xrightarrow{\text{cod}} \mathcal{C}$$

$\text{Rel}_{(E,M)}(\mathcal{C})$

The guiding question :

What makes a double category of relations ?



Rel , $\text{Rel}(\mathcal{C})$, Span ,
 $\text{Rel}_{(E,M)}(\mathcal{C})$, $\mathcal{Q}\text{-Rel}$, ...



Prof , Par , ...

My answer : the Frobenius axiom

1. Background
2. The Frobenius axiom
3. Bonus : Maps revisited

1. Background

2. The Frobenius axiom

3. Bonus : Maps revisited

Bicategories of relations

03 / 17

Allegory

[Freyd, Scedrov, '90]

Formulation in terms of **involution**.

$$(-)^{\circ} : B(I, J) \longrightarrow B(J, I)$$

Cartesian bicategory

[Carboni, Walters '87, Carboni, et.al. '07,
Todd, nLab post.]

Formulation using **adjoints (= maps)**
in bicategories.

Theorem.

B is a unital
tabular allegory

$\iff$

$B \simeq \text{Rel}(\mathcal{C})$
for a regular category $\mathcal{C}$

$\iff$

B is a functionally-
complete discrete
locally-posetal
cartesian bicategory.

Bicategories of relations, generalized

04/17

Definition. $p : \mathcal{E} \rightarrow \mathcal{B}$: fibration, $\mathcal{B}$ has finite products, $I, J \in \mathcal{B}$

| A p -relation $I \rightarrowtail J$ is an object $R \in \mathcal{E}_{I \times J}$. the fiber over $I \times J$

Proposition. [Pavlovic '96, Lawler '15]

| $p : \mathcal{E} \rightarrow \mathcal{B}$: a regular fibration $\rightsquigarrow R_p$: a bicategory of p -relations.

Existence of a (partial) right adjoint is stated without proof.

Theorem. [Bonchi, Santamaria, Seeber, Sobocinski '21]

| $\left(\begin{array}{c} \text{elementary existential} \\ \text{doctrines} \end{array} \right) \xrightleftharpoons[\perp]{R-} \left(\begin{array}{c} \text{'bicategories of} \\ \text{'relations' in [CW87]} \end{array} \right)$

Double categories

Definition.

A (pseudo) double category consists of

- objects $A, B, \dots$

- tight arrows

$$\begin{array}{c} A \\ \downarrow f \\ B \end{array}, \dots$$

- loose arrows

$$A \xrightarrow{p} B, \dots$$

- cells

$$\begin{array}{ccccc} A & \xrightarrow{p} & B & & \\ f \downarrow & \sigma & \downarrow g & & \\ C & \xrightarrow{q} & D & & \end{array}, \dots$$

with compositions of
tight arrows,

loose arrows,

$$A \xrightarrow{p} K \xrightarrow{m} X$$

and cells

$$\begin{array}{ccccc} A & \xrightarrow{p} & K & \xrightarrow{m} & X \\ f \downarrow & \sigma & \downarrow s & \nu & \downarrow u \\ B & \xrightarrow{q} & L & \xrightarrow{n} & Y \end{array}$$

$$\begin{array}{c} A \\ \downarrow f \\ B \\ \downarrow g \\ C \end{array}$$

$$\begin{array}{ccccc} A & \xrightarrow{p} & K & & \\ f \downarrow & \sigma & \downarrow s & & \\ B & \xrightarrow{q} & L & & \\ g \downarrow & \tau & \downarrow t & & \\ C & \xrightarrow{r} & M & & \end{array}$$

subject to some coherence conditions.

Double categories

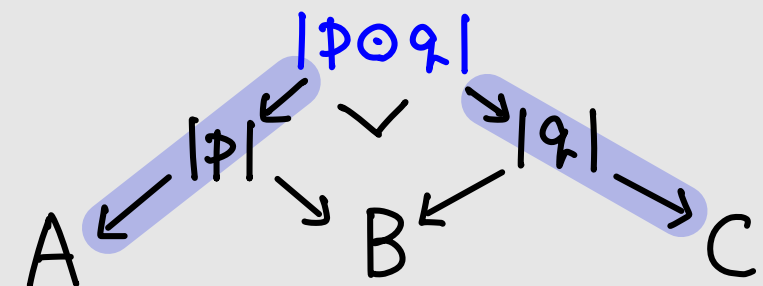
06 / 17

<u>Example</u>	Rel	Span	Prof
Objects	sets	sets	categories
Tight arrows	functions	functions	functors
Loose arrows	relations	spans	profunctors
Cells	$ \begin{array}{ccc} A & \xrightarrow{p} & B \\ f \downarrow & \tau & \downarrow g \\ C & \xrightarrow{q} & D \end{array} $ A cell exists iff $ \begin{array}{c} p(a, b) \\ \Downarrow \\ q(f(a), g(b)) \end{array} (\forall a, b) $	$ \begin{array}{ccc} A & \xleftarrow{ p } & B \\ f \downarrow & \eta & \downarrow h & \eta & \downarrow g \\ C & \xleftarrow{ q } & D \end{array} $	$ \begin{array}{ccc} A^{\text{op}} \times B & \xrightarrow{p} & \text{Set} \\ f^{\text{op}} \times g \downarrow & \Downarrow \tau & \\ C^{\text{op}} \times D & \xrightarrow{q} & \end{array} $

Remark

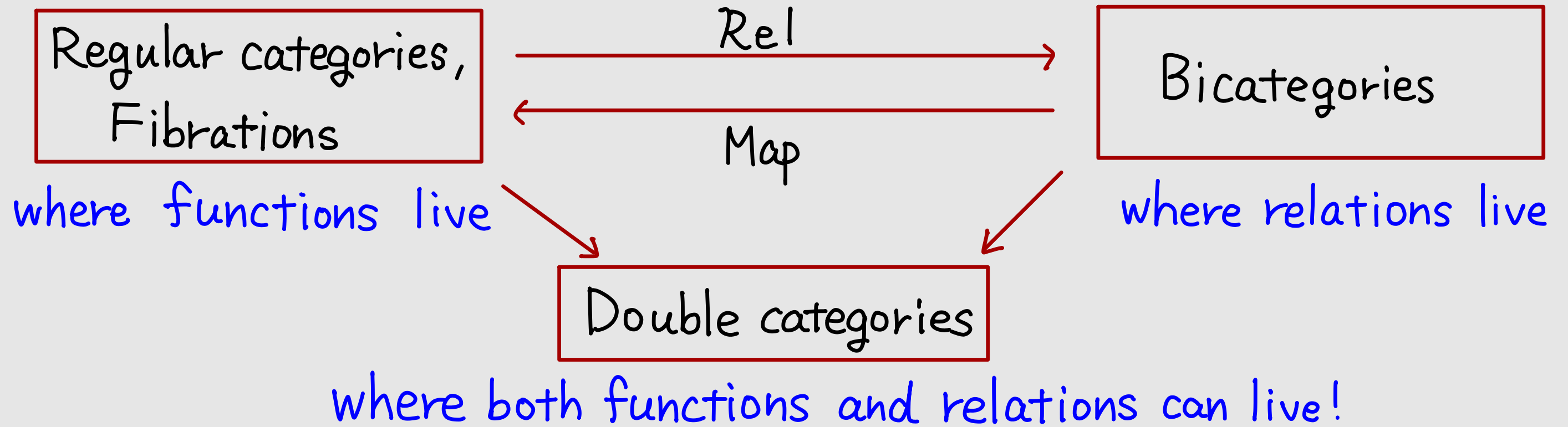
The loose composition is required to be associative *only up to isomorphism*.

e.g. the composition of spans



Motivating double categories

07 / 17



- ① can express the interaction b/w functions & relations with companions & conjoints.
- ② can distinguish functions from maps ($\doteq$ functional relations)

Motivating double categories

③ Cartesian double categories are easier to handle.

Definition. $\mathbb{D}$: double category

$\mathbb{D}$ is **cartesian** if $\mathbb{D} \begin{array}{c} \xrightarrow{!} \\ \perp \\ \xleftarrow{E} \end{array} \mathbb{1} \quad \& \quad \mathbb{D} \begin{array}{c} \xrightarrow{\Delta} \\ \perp \\ \xleftarrow{E \times} \end{array} \mathbb{D} \times \mathbb{D} \quad \text{in } \mathbf{Dbl}.$

Proposition. [Patterson '24, N.]

$\mathbb{D}$: cartesian equipment $\implies \mathcal{L}(\mathbb{D})$: a cartesian bicategory

Remark Most of known examples of cartesian bicategories arise in this way. For a cartesian bicategory B , there is a cartesian double **bicategory** whose loose part is B . [Verity'92]

1. Background

2. The Frobenius axiom

3. Bonus : Maps revisited

The Frobenius axiom

09 / 17

Definition 2.1. (i) An object X in a Cartesian bicategory is discrete when the multiplication Δ_X^* and the comultiplication Δ_X satisfy

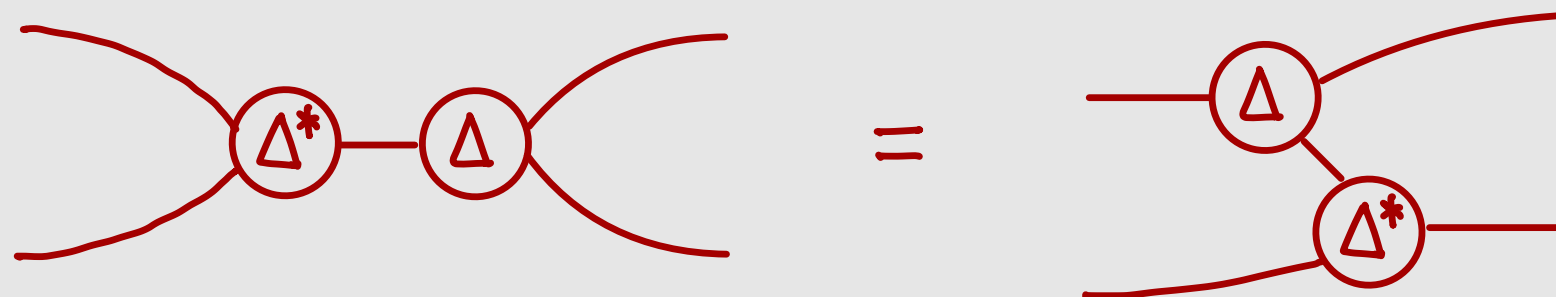
$$(D) \quad \Delta \cdot \Delta^* = (\Delta^* \otimes 1) \cdot (1 \otimes \Delta).$$

(We are forgetting the associativity in the middle.)

(ii) A Cartesian bicategory is called a 'bicategory of relations' if every object is discrete.

[Carboni, Walters. '87]

This condition was later rephrased as the Frobenius axiom, because this is equivalent to the condition for Frobenius algebra. [Walters, Wood '08]



The Frobenius axiom

Definition. [Walters, Wood '08, Lambert '23, Hoshino, N. '25, N.]

$\mathbb{D}$: cartesian equipment

$I \in \mathbb{D}$ is Frobenius if

$$\begin{array}{c}
 \begin{array}{ccc}
 & I & \\
 \Delta \swarrow & & \searrow \Delta \\
 I \times I & \xrightarrow{\text{opc.}} & I \times I \\
 \Delta \times \text{id} \searrow & & \swarrow \text{id} \times \Delta \\
 & I \times I \times I &
 \end{array}
 = \begin{array}{c} \bullet \\ \swarrow \quad \searrow \\ \bullet = \bullet \\ \swarrow \quad \searrow \\ \bullet \end{array}
 \Leftrightarrow \left(\begin{array}{ccccc}
 I \times I & \xrightarrow{\Delta^*} & I & \xrightarrow{\Delta^*} & I \times I \\
 \parallel & & \text{SII} & & \parallel \\
 I \times I & \xrightarrow{(\Delta \times \text{id})_*} & I \times I \times I & \xrightarrow{(\text{id} \times \Delta)^*} & I \times I
 \end{array} \right)
 \end{array}$$

canonically.

$\mathbb{D}$ is Frobenius if every object in $\mathbb{D}$ is Frobenius.

Example. $\mathcal{C}$: a regular category $\Rightarrow \text{Rel}(\mathcal{C})$ is Frobenius

For $A \in \mathcal{C}$ and $(x, y) : A \times A$, $(z, w) : A \times A$,

$$\exists u : A \ ((x, y) = (u, u) \wedge (u, u) = (z, w)) \iff (x, x, y) = (z, w, w)$$

Example. In Prof, $\mathcal{C}$ is Frobenius $\Leftrightarrow \mathcal{C}$ is a groupoid.

11 / 17

Proof sketch [Walters, Wood '08]

$I \in \mathbb{D}$ is Frobenius if $\Delta^* \Delta_* \cong (\Delta \times \text{id})_* (\text{id} \times \Delta)^*$

$$(\Delta^* \Delta_*)((c_1, c_2), (d_1, d_2)) = \int^{e \in \mathcal{C}} \mathcal{C}(c_1, e) \times \mathcal{C}(c_2, e) \times \mathcal{C}(e, d_1) \times \mathcal{C}(e, d_2)$$

$$\downarrow$$

$$((\Delta \times \text{id})_* (\text{id} \times \Delta)^*)((c_1, c_2), (d_1, d_2)) = \mathcal{C}(c_1, d_1) \times \mathcal{C}(c_1, d_2) \times \mathcal{C}(c_2, d_2)$$

$$\left[\begin{array}{ccc} c_1 & \xrightarrow{\text{red}} & e \\ & \searrow \text{blue} & \\ c_2 & \xrightarrow{\text{green}} & e \end{array} \begin{array}{ccc} e & \xrightarrow{\text{red}} & d_1 \\ & \searrow \text{blue} & \\ e & \xrightarrow{\text{green}} & d_2 \end{array} \right] \mapsto \left(\begin{array}{ccc} c_1 & \xrightarrow{\text{red}} & d_1 \\ & \searrow \text{blue} & \\ c_2 & \xrightarrow{\text{green}} & d_2 \end{array} \right)$$

$(\Rightarrow)$ Take $f: a \rightarrow b$ in $\mathcal{C}$.

$$\exists \begin{array}{ccccc} a & & f_1 & & a \\ & \searrow & & \nearrow & \\ & e & & & \\ & \nearrow & & \searrow & \\ b & & f_2 & & b \end{array}$$

$$\mapsto \begin{array}{ccc} a & = & a \\ & \searrow f & \\ b & = & b \end{array}$$

$b \xrightarrow{f_2} e \xrightarrow{f_3} a$ is
the inverse of f .

$(\Leftarrow)$ A small puzzle! $\square$

The Frobenius axiom induces involution

12 / 17

Slogan : The Frobenius axiom is an axiom of symmetry.

Proposition. [Hoshino, N. '25, N.]

$\mathbb{D}$: a Frobenius cartesian equipment.

Every A and $\downarrow f$ in $\mathbb{D}$ are **self-dual** in the loose direction :

$$\mathbb{D}_\ell(C \times A, B) \cong \mathbb{D}_\ell(C, A \times B), \quad \mathbb{D}_c(h \times f, g) \cong \mathbb{D}_c(h, f \times g)$$

The unit and counit on A are

$$1 \xrightarrow{!^*} A \xrightarrow{\Delta^*} A \times A \quad \text{and} \quad A \times A \xrightarrow{\Delta^*} A \xrightarrow{!^*} 1.$$

Cor. $\mathcal{L}(\mathbb{D})$ is a compact closed bicategory in the sense of [Stay '16].

From equipments to fibrations

Definition. $\mathbb{D}$: cartesian equipment $\rightsquigarrow$

a category $\mathcal{U}_{\mathbb{D}}$: objects : $A \xrightarrow{\alpha} 1$,

morphisms $(A \xrightarrow{\alpha} 1) \rightarrow (B \xrightarrow{\beta} 1) : \begin{array}{ccc} A & \xrightarrow{\alpha} & 1 \\ f \downarrow & \varphi & \parallel \\ B & \xrightarrow[\beta]{} & 1 \end{array}$

a fibration $\mathcal{U}_{\mathbb{D}} : \mathcal{U}_{\mathbb{D}} \longrightarrow \mathbb{D}_0 ; (A \xrightarrow{\alpha} 1) \mapsto A$

When $\mathbb{D}$ is Frobenius, all information in $\mathbb{D}$ is recoverable from $\mathcal{U}_{\mathbb{D}}$.

$$\begin{array}{ccc} \begin{array}{ccc} C & \xrightarrow{\alpha} & A \\ f \downarrow & \tau & \downarrow g \\ D & \xrightarrow[\beta]{} & B \end{array} & \rightsquigarrow & \begin{array}{ccc} C \times A & \xrightarrow{\tilde{\alpha}} & 1 \\ f \times g \downarrow & \tilde{\tau} & \parallel \\ D \times B & \xrightarrow[\tilde{\beta}]{} & 1 \end{array} & \rightsquigarrow & \tilde{\alpha} \xrightarrow{\tilde{\tau}} \tilde{\beta} \text{ in } \mathcal{U}_{\mathbb{D}} \end{array}$$

Proposition. [N.] $\mathbb{D}$: Frobenius $\Rightarrow$

$\mathcal{U}_{\mathbb{D}} : \mathcal{U}_{\mathbb{D}} \longrightarrow \mathbb{D}_0$ is an elementary existential fibration.

Main results

Pavlovic's bicategories of predicates are extended to double categories.
[Lawler 15, N.]

Elementary existential fibrations $\xrightarrow{\text{Bil}}$ Cartesian equipment

$\text{Bil}(\mathbb{P})$: the double category of $\mathbb{P}$ -relations

Theorem. [N.]

$\mathbb{P}$: an e.e. fibration $\rightsquigarrow \text{Bil}(\mathbb{P})$: a Frobenius cartesian equipment.

This extends to a biequivalence:

$$\mathcal{E}\mathcal{E}\text{Fib} \xrightleftharpoons[\cong]{\text{Bil}} \text{CartEq}_{\text{Frob}}$$

$\mathbb{D} \simeq \text{Bil}(\mathbb{P})$ for some $\mathbb{P}$ iff it is Frobenius.

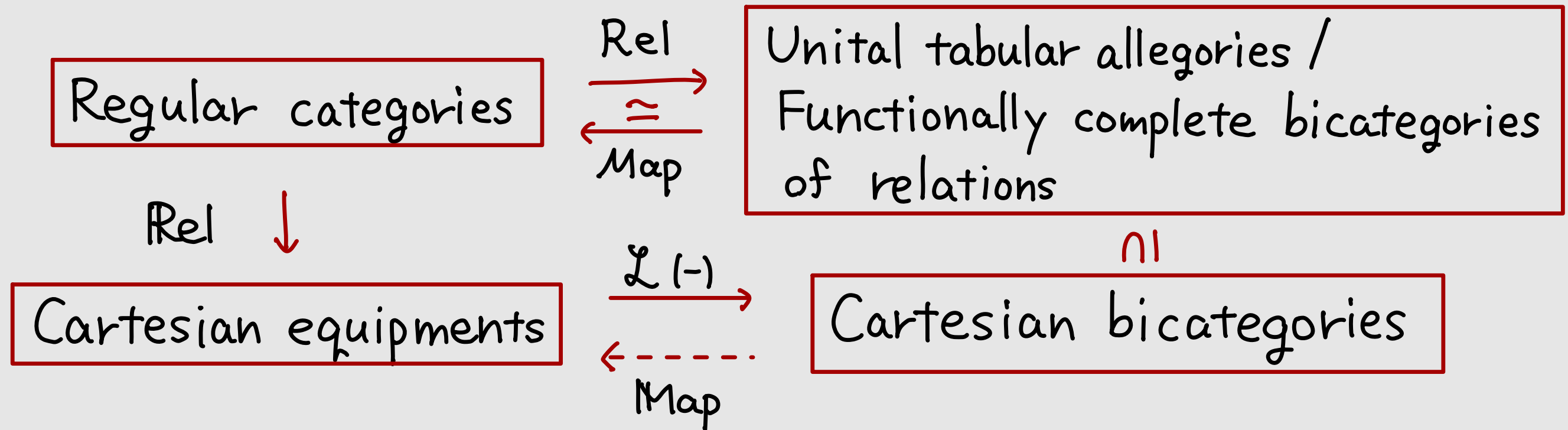
Example. $\text{Prof}_{\text{grpd}} \simeq \text{Bil}(\text{GrpdAct} \rightarrow \text{Grpd})$

1. Background

2. The Frobenius axiom

3. Bonus : Maps revisited

Definition. A **map** in a bicategory is a left adjoint 1-cell in it.



Definition. A bicategory $\mathcal{B}$ is **map discrete** if $\text{Map}(\mathcal{B})(I, J) \subseteq \mathcal{B}(I, J)$
is equivalent to discrete categories.

This is sufficient to construct a double category $\mathbb{M}ap(B)$.

Revisiting Map

Definition. [Paré 21] An equipment $\mathbb{D}$ is **Cauchy** if every map
 | is isomorphic to f_* for some $f: A \rightarrow B$.

Unit-
pure

Map discrete
Cauchy cart. equipments

$\mathcal{H}(-)$
 $\xrightarrow{\quad \simeq \quad}$
 $\text{Map}(-)$

Map discrete
cart. bicategories

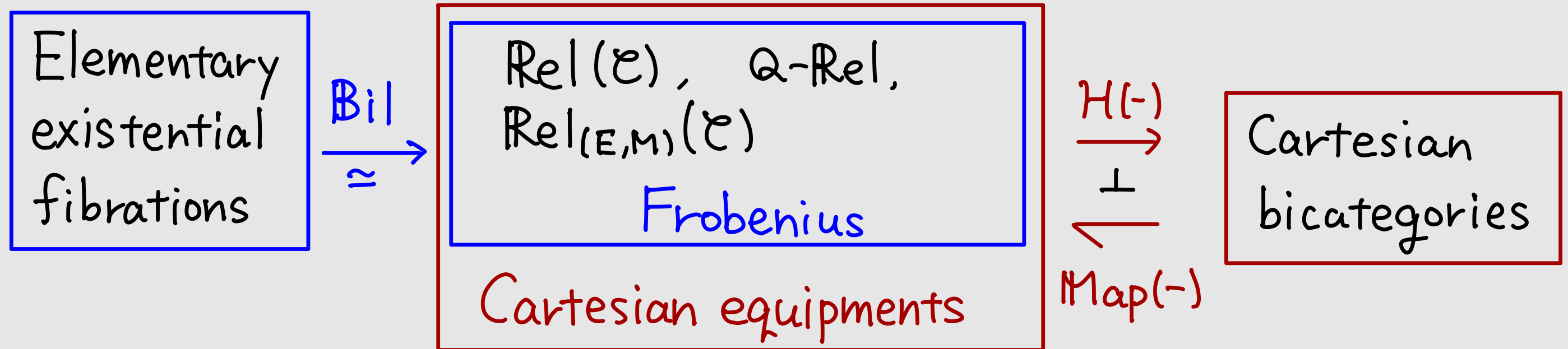
Proposition. [N.]

$\text{CartEq}_{\text{Cau}} \xleftarrow{\text{Cau}} \text{CartEq}_{\text{MD}} \downarrow$
 $\text{CartEq}_{\text{Cau}} \xrightarrow{\perp} \text{CartEq}$

Corollary. [Bonchi, et al. '21, N.] [Pavlovic, '96]

$\text{EEFib} \xrightarrow{\text{locpos}} \text{CartEq}_{\text{Frob, MD}} \xleftarrow{\text{Cau}} \text{CartEq}_{\text{Cau, Frob, MD}} \xrightarrow{\text{locpos}} \text{CartBi}_{\text{Frob, MD}} \xleftarrow{\text{Map}(-)} \text{Map}(-)$

Summary



Future Work

- Generalizing to monoidal equipments.
- Studying categorical logic with cartesian equipments (quotient completion, tabulators)

Reference

- E. Aleiferi. Cartesian Double Categories with an Emphasis on Characterizing Spans. PhD thesis, Dalhousie Univ, 2018.
- A. Carboni & R.F.C. Walters. Cartesian bicategories I. JPAA, 49(1-2): 11-32, 1987.
- F. Bonchi, A. Santamaria, J. Seeber, P. Sobociński. On Doctrines and Cartesian Bicategories. CALCO2021, LIPIcs 10:1-10:17.
- K. Hoshino & H. Nasu. Double categories of relations relative to factorisation system. (ACS, 2025.)
- E. Patterson. Transposing cartesian and other structure in double categories. preprint, 2024
- D. Pavlović. Maps II. Chasing diagrams in categorical proof theory. J.IGPL, 4(2): 159-194, 1996
- M. Shulman. Framed bicategories and monoidal fibrations TAC. 20, No.18. 2008.
- M. Stay. Compact closed bicategories. TAC. 28. 616-696. 2013.
- T. Trimble and others. cartesian bicategory. nLab. 2009. version 27
- D. Verity. Enriched categories, internal categories and change of base. PhD thesis, 1992. (Repr. TAC. 36: No.8. 206-249, 2011.)
- R.F.C. Walters, R.J. Wood. Frobenius objects in Cartesian bicategories. TAC. 20. No.3: 25-47. 2008.

Thank you!

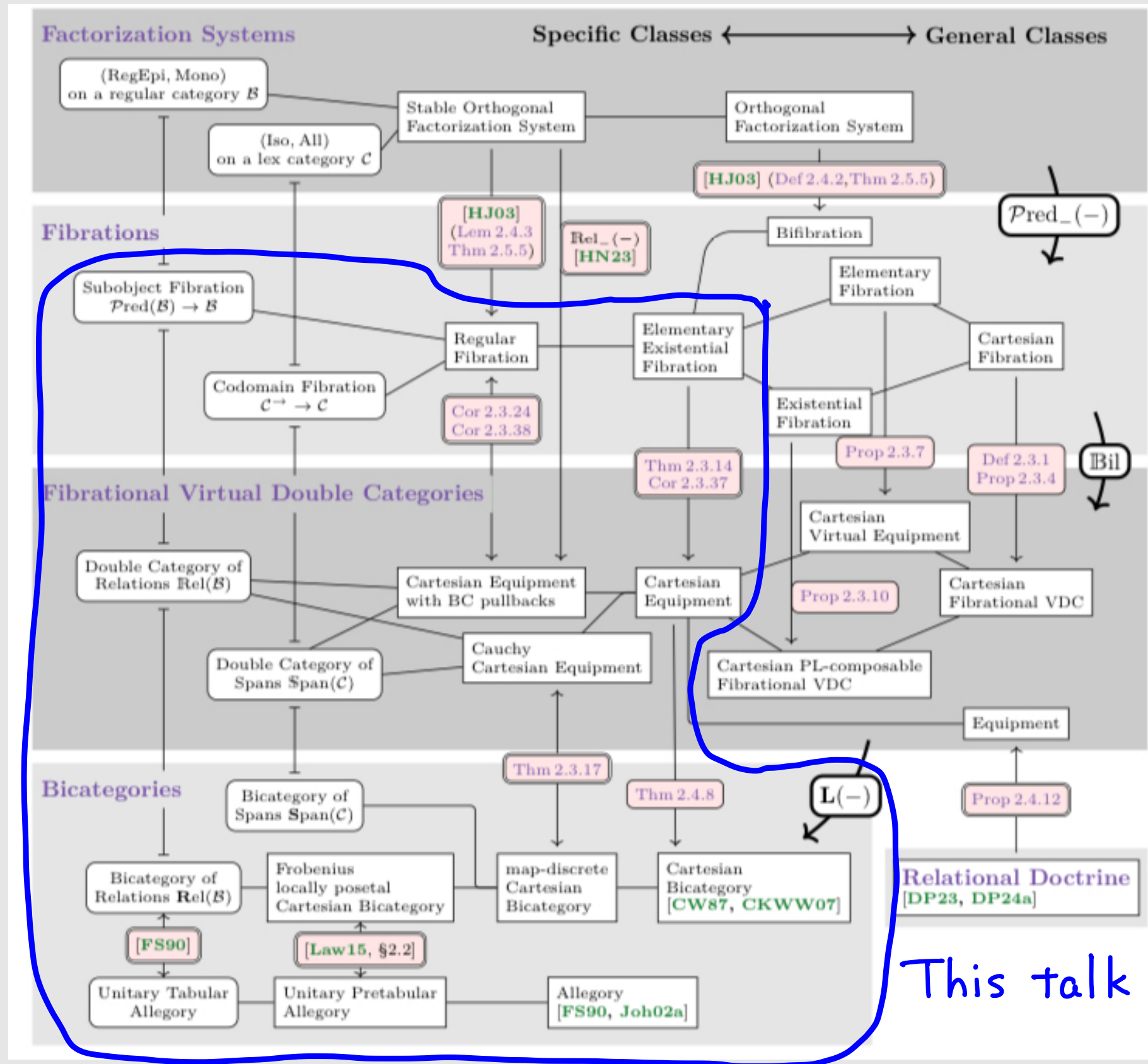
my homepage :

hayatonasu.github.io

my thesis :

[arXiv:2501.17869](https://arxiv.org/abs/2501.17869)

talk at CMS:



Characterization Results

A1 / A3

Theorem. [Lambert '23, Hoshino.N. '25]

$\mathbb{D} \simeq \text{Rel}(\mathcal{C})$ for some regular category

$\Leftrightarrow \begin{cases} \mathbb{D} : \text{locally posetal discrete cartesian equipment} \\ \text{with Mono-comprehension scheme} \end{cases}$

Theorem. [Hoshino.N. '25, my thesis]

$\mathbb{D} \simeq \text{Rel}_{(E,M)}(\mathcal{C})$ for some SOFS (E,M)

$\Leftrightarrow \begin{cases} \mathbb{D} : \text{cartesian equipment with BC pullbacks \& effective} \\ \text{tabulators} \\ \text{and } \text{Fib}(\mathbb{D}) \text{ is closed under composition.} \end{cases}$

Bicategories of relations

A2 / A3

- 1984 • "Bicategories of spans and relations" Carboni, Kasangian, Street.
defined $\text{Span}(\mathcal{C})$ & $\text{Rel}(\mathcal{C})$
- 1987 • "Cartesian bicategories I" Carboni, Walters.
- 1990 • "Categories, Allegories" Freyd, Scedrov.
- 1996 • "Maps II: Chasing diagrams in categorical proof theory" Pavlovic
defined bicategories of predicates for regular fibrations
- 2021 • "On Doctrines and Cartesian Bicategories" Bonchi et.al.

Double categories of relations and spans

A3 / A3

- 1999 + "Limits in double categories". Grandis, Paré.
- 2012 + "Span, cospan, and other double categories" Niefield.
- 2019 + "Cartesian Double categories with an Emphasis on Characterizing Spans" Aleiferi. characterized $\text{Span}(\mathcal{C})$
- 2023 + "Double categories of relations" Lambert
characterized $\text{Rel}(\mathcal{C})$ for regular categories $\mathcal{C}$
- 2025 + "Double categories of relations relative to factorization systems" Hoshino, N. characterized $\text{Rel}_{(E,M)}(\mathcal{C})$.